A Modular Model Checker for Reference Nets: MoMoC

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m MoMoC}$ pursues two goals:

- Teaching model checking
- Extensibility, to form a basis for further research on model checking of Reference nets

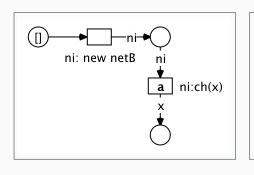
Reference Nets

The Java Reference Net Formalism

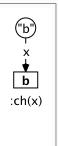
- Coloured Petri net (CPN) formalism
- ullet Primary formalism of the Renew simulator
- Java code inscriptions
- Tokens: Java objects or net instances
- Interaction: synchronous channels

Net Instance Tokens

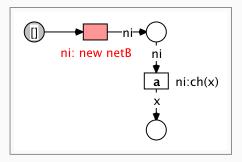
- nets-within-nets
- follows the object-oriented paradigm
 - net template class
 - **net instance** instance/object
 - **net elements** internals of a class
 - uplinks of a net methods/interface of a class
 - invoking an uplink invoking a method

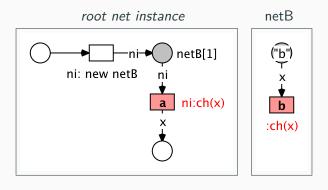




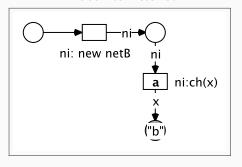


root net instance





root net instance



Features

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Modular Model Checker (MoMoC)

- Explicit CTL-Model Checking for Reference Nets
- Parsing
- Result visualization (exploration, colorization, layouting)
- Net Instance Quantifier
- Simpler net formalisms (P/T nets, CPNs) can be treated as flat Reference nets

Atomic Propositions

- FIREABLE(T)
- DEADLOCK

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- Marking predicates..

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- $?(\textit{Net}, p) \equiv$ There exists a net instance of the template Net that satisfies p.

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Scales independently of the size of the reachability graph, however net instances cannot be tracked over multiple states.

Parsing

Uses ANTLR as a framework for parsing.

Parsing features of MoMoC:

- Parsing of different notations
- Normalization
- Reduction
- Encoding

Result Visualization

Goal: Comprehensive results that help teaching (CTL) Model Checking

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- States of the RG can be explored
- RG can interactively be colorized with results of subroutines

Architecture

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 Binding Core - Finds bindings and calculates successive markings, thus defines the semantics

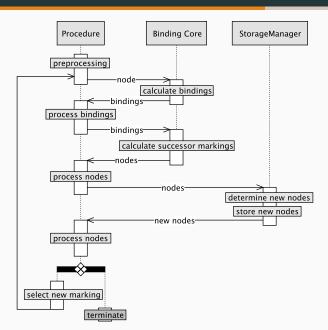
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- Storage Manager Stores the reachability graph and finds cycles in the graph

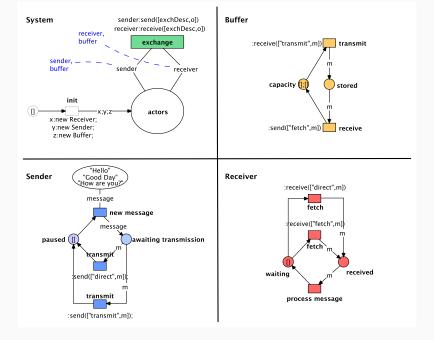
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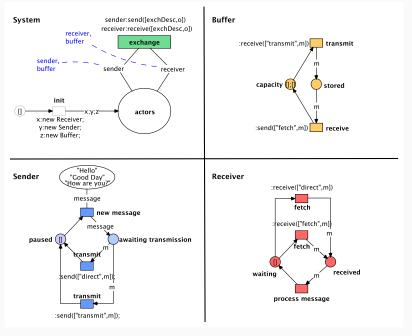
- Binding Core Finds bindings and calculates successive markings, thus defines the semantics
- Storage Manager Stores the reachability graph and finds cycles in the graph
- Procedure Contains logic and steps to process a query

Module Interaction



Demo





 $EG(!(Buffer, m(stored) = 0) \land AF?(Receiver, m(received) = 1))$

• • • EG[AND[![Buffer, m[stored]=0],NOT[EG[NOT[?[Receiver, m[received]=1]]]]]] ▼ AND[![Buffer, m[stored]=0].NOT[EG[NOT[?[Receiver, m[received]=1]]]]] Specification ![Buffer, m[stored]=0]

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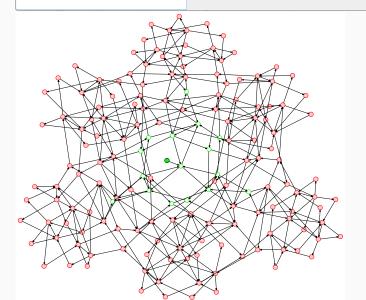
Result

EG(!(Buffer,m(stored)=0) ^ AF?(Receiver,m(received)=1))

Result: TRUE Root Net: System[10786]

Normalized Spec.:EG[AND[![Buffer, m[stored]=0],NOT[EG[NOT[?[Receiver, m[received]=1]]]]]]

Show Reachability Graph



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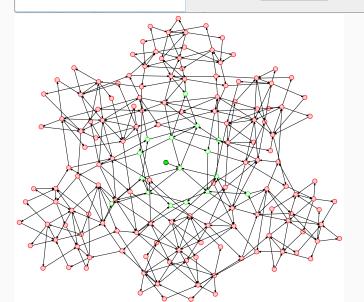
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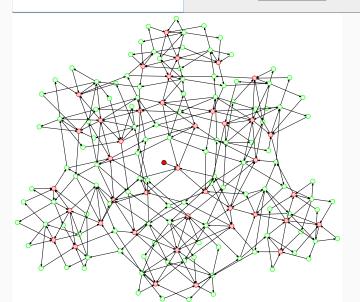
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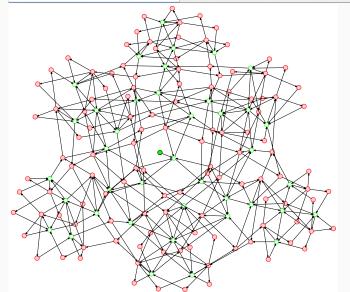
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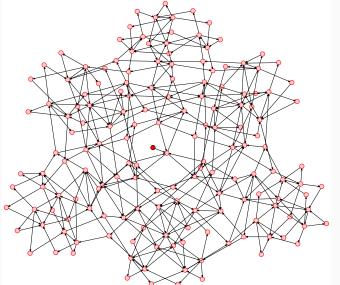
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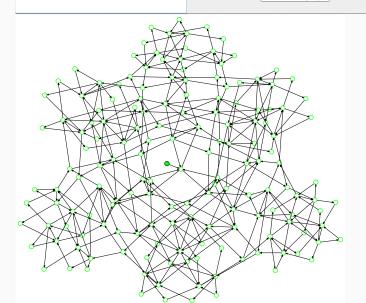
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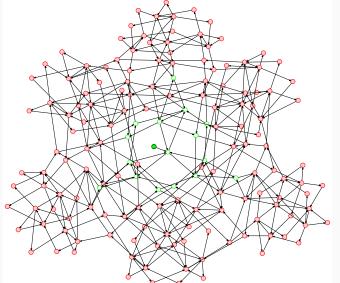
| Result
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Show Reachability Graph

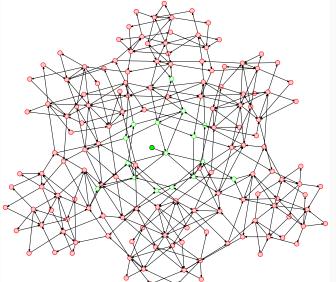
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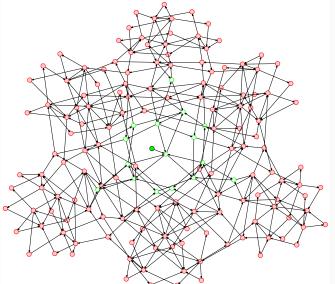


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Evaluation

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- \bullet Teaching-size problems ($<\!10k$ states) are unproblematic with average computing power
- Colorization is helpful

Conclusion

- CTL Model-Checking of Reference nets
- Modular architecture
- Net Instance Quantifier
- Result visualization

Outlook

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Teaching-oriented goals:

- LTL-Model Checking
- Coverability graph
- Visualization of large RGs
 - \rightarrow Interactive trace visualization

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- LTL-Model Checking
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Efficiency-oriented goals:

- Code-specific improvements
- Transfer of known techniques to Reference nets
- Techniques that exploit the structural information contained in Reference nets